

Shape Interrogation II

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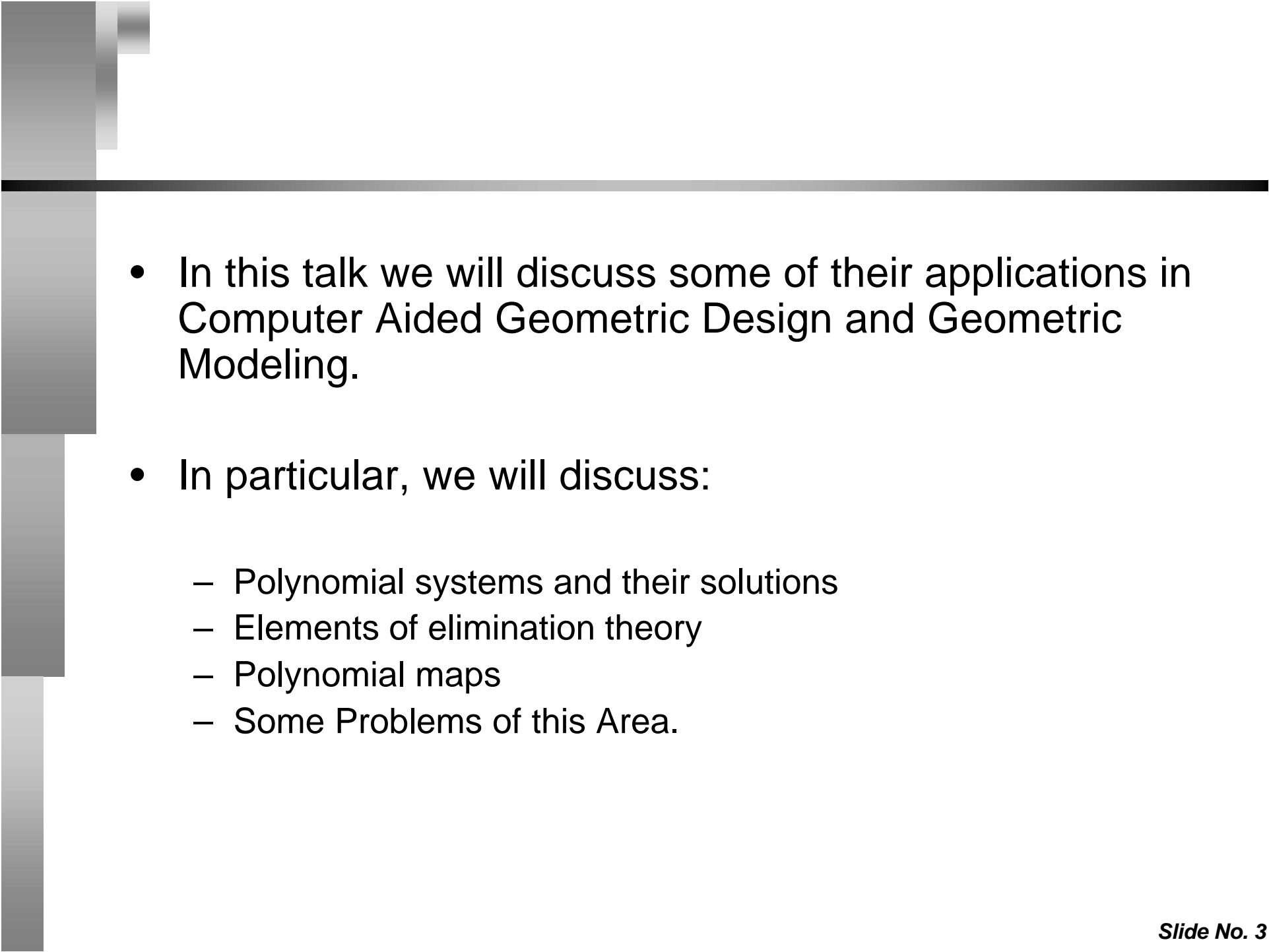
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Introduction

- Polynomials are used in various branches of computational science.
- They can be found in mathematics, computer science, engineering and many other fields.
- There are two basic reasons for that:
 - Most functions can be approximated by polynomial functions, and
 - They are rather easy to use in a computer code.
- Thus, they serve as good substitutes for functions that are difficult to deal with.

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- In this talk we will discuss some of their applications in Computer Aided Geometric Design and Geometric Modeling.
 - In particular, we will discuss:
 - Polynomial systems and their solutions
 - Elements of elimination theory
 - Polynomial maps
 - Some Problems of this Area.

A Strange Example

- As an indication of the difference in moving from one dimension to the next, even for simple functions—like polynomials—let us consider the following:
- Example 1. Every polynomial function $y = p(x)$ with $p(x) > 0, \forall x \in \mathbb{R}$ has at least one (real) critical point.

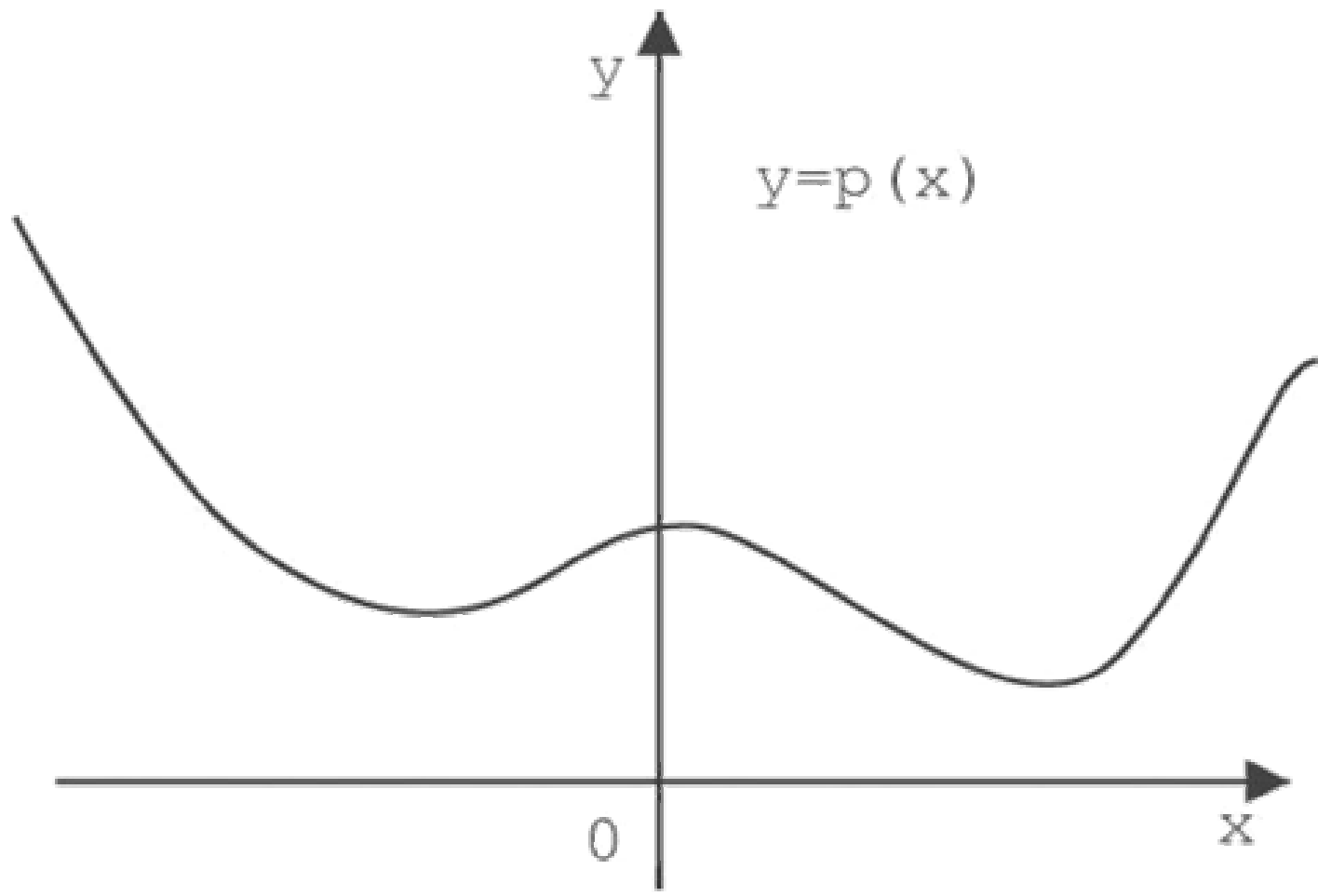
$$\left[\lim_{|x| \rightarrow \infty} p(x) = \infty. \right]$$

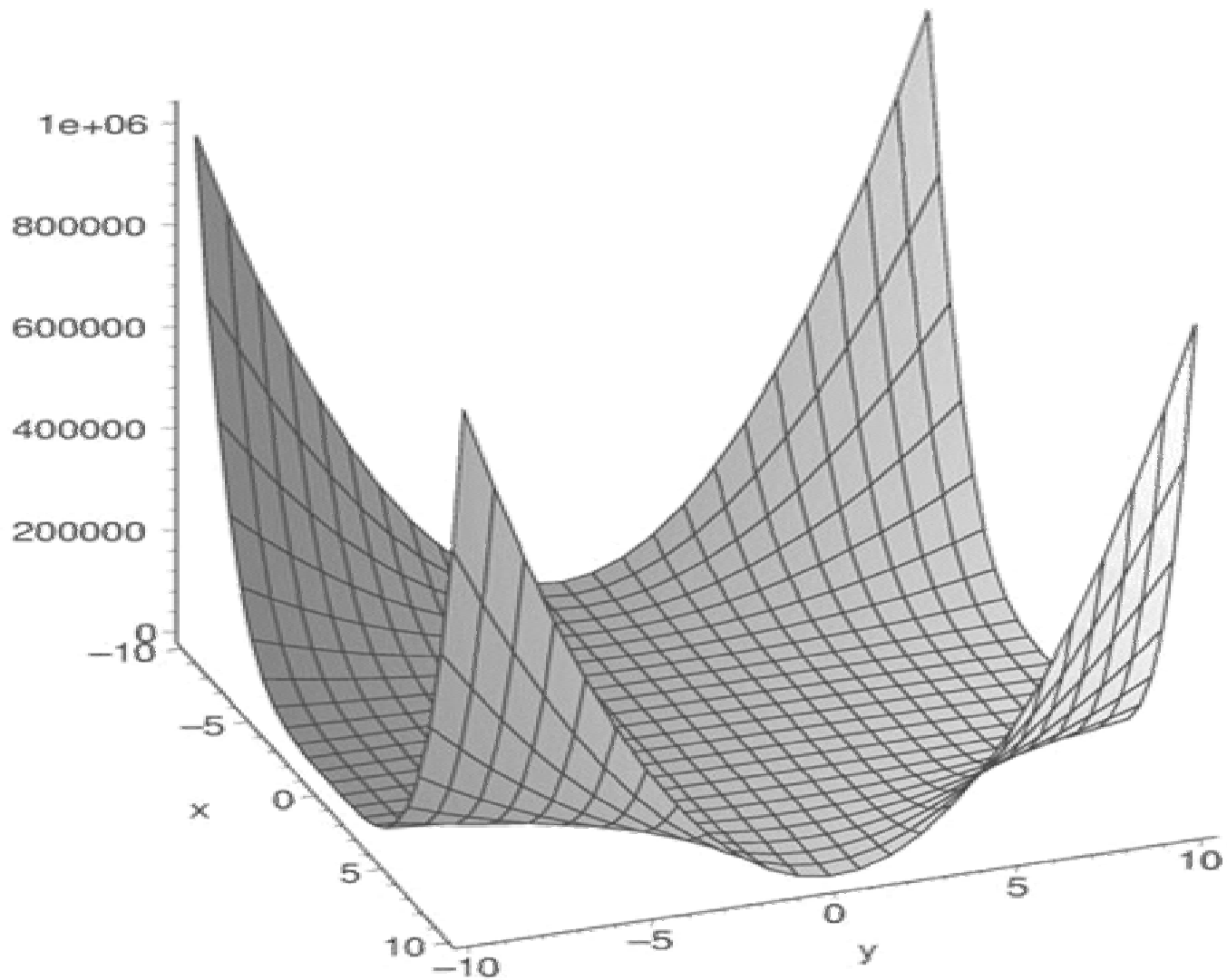
- Example 2. The polynomial

$$p(x, y) = (x^2 y - x - 1)^2 + x^2$$

- has the property that, for every $(x, y) \in \mathbb{R}^2, p(x, y) > 0,$
- but the function $p(x, y)$ does not have any (real) critical point.

$$\left[\lim_{|(x, y)| \rightarrow \infty} p(x, y) \text{ Does not exist.} \right]$$





Polynomial Systems

- Polynomials are popular in curve and surface representation.
- Thus, many critical problems in CAGD, such as surface interrogation, are reduced to finding the zero set of a system of polynomial equations

$$f(x) = 0$$

where $f = (f_1, \dots, f_n)$ and each f_i is a polynomial of m independent variables $x = (x_1, \dots, x_m)$.

Polynomial Systems

- Several root-finding methods for polynomial systems have been used in practice.
- These can be categorized as:
 - Algebraic and hybrid methods,
 - Homotopy methods, and
 - Subdivision methods.
- Among those types, the subdivision methods have been widely used in practice.
- The Interval Projected Polyhedral (IPP) algorithm is one example, and it has successfully been applied to various problems.

Motivation

- **Difficulties in handling roots with high multiplicity**
 - Performance deterioration
 - Lack of robustness in numerical computation
 - Round-off errors during floating point arithmetic
- **Limited research on root multiplicity of a system of equations**
 - Heuristic approaches are needed for practical purposes.

Objectives

- **Develop practical algorithms to isolate and compute roots and their multiplicities.**
- **Improve the Interval Projected Polyhedron (IPP) algorithms.**

Multiplicity of Roots

- **Univariate Case**

- A root a of $f(x)=0$ has multiplicity k if

$$f(a) = f'(a) = \cdots = f^{(k-1)}(a) = 0, \text{ and } f^{(k)}(a) \neq 0$$

- **Bivariate Case**

- Define

$$V_f = \{(x, y) \in \mathbf{C} \mid f(x, y) = 0\}$$

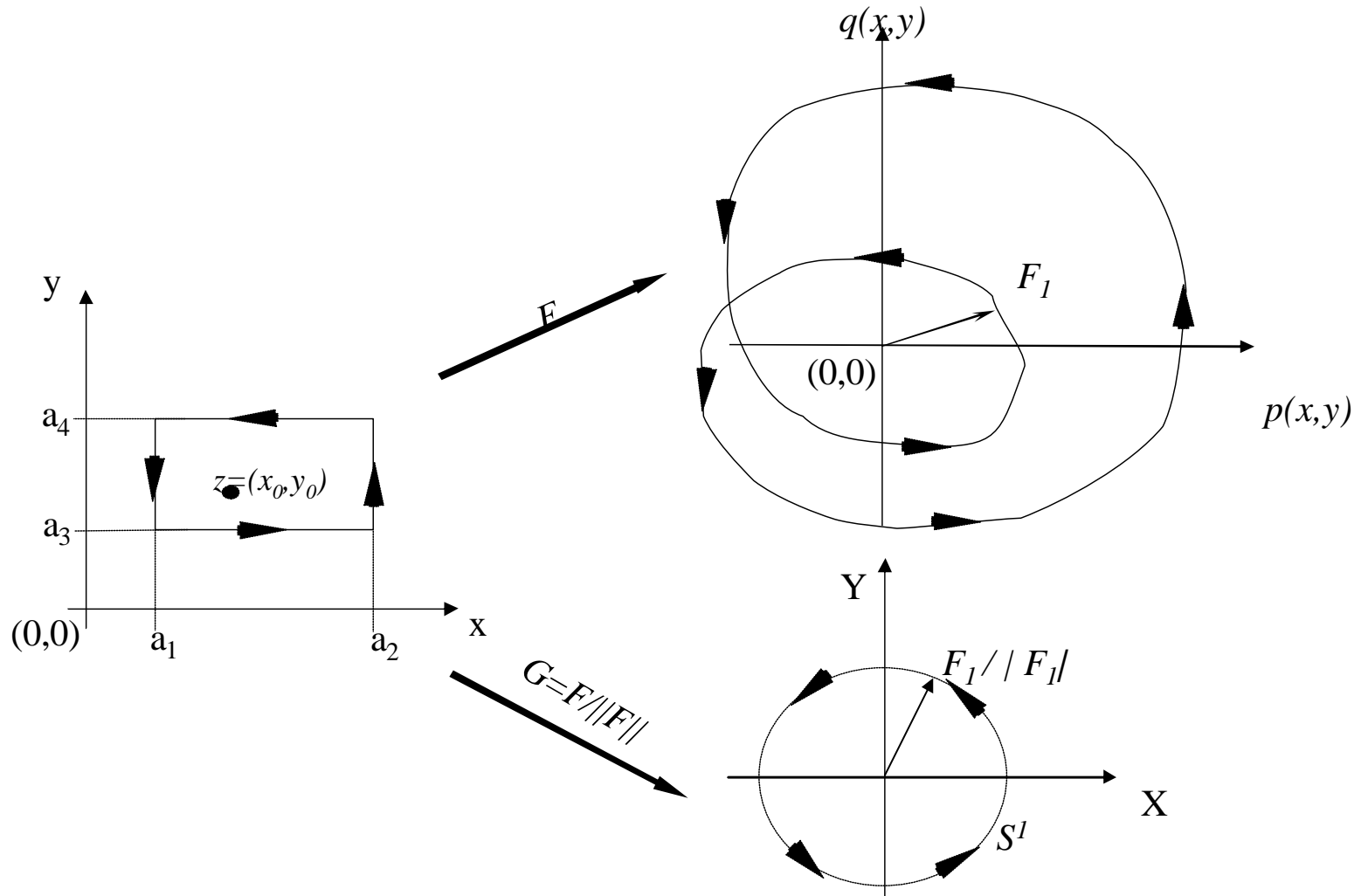
$$V_g = \{(x, y) \in \mathbf{C} \mid g(x, y) = 0\}$$

- Suppose that z_0 is the only common point of V_f and V_g lying above x_0 . Consider $h(x)=\text{Res}_y(f,g)$, the resultant of f,g with respect to y . Then the multiplicity of $z_0=(x_0,y_0)$ as a root of the system is the multiplicity of x_0 as a zero of $h(x)$.

Degree of the Gauss Map

- Let $p(x,y)$, $q(x,y)$ be polynomials with rational coefficients without common factors, of degrees n_1 and n_2 , and let $F=(p, q)$.
- Let A be a rectangle in the plane defined by $a_1 \leq x \leq a_2$, $a_3 \leq y \leq a_4$, $a_1 < a_2$, $a_3 < a_4$, $a_i \in \mathbf{Q}$, $i=1,2,3,4$ so that no zero of F lies its boundary ∂A , and $p \cdot q$ does not vanish at its vertices.
 - Gauss map $G:\partial A \rightarrow S^1$, $G = F/\|F\|$, where S^1 is the unit circle.
 - G is continuous ($\|F\| \neq 0$ on ∂A).
 - ∂A and S^1 carry the counterclockwise orientation.
- Degree d of G : an integer indicating how many times ∂A is wrapped around S^1 by G .

Illustration of the Gauss Map



The Cauchy Index

- **Preliminaries**

- $R(x)$: a rational function $q(x)/p(x)$, where p, q are polynomials.
- $[a,b]$: a closed interval, $a < b$. R does not become infinite at the end points.

- **Definition of the *Cauchy index***

By the *Cauchy index*, $I_a^b R$ of R over $[a,b]$, we mean $I_a^b R = N_-^+ - N_+^-$ where $N_-^+ (N_+^-)$ denotes the number of points in (a,b) at which $R(x)$ jumps from $-\infty$ to $+\infty$ ($+\infty$ to $-\infty$), respectively, as x is moving from a to b . Notice that $I_a^b R = -I_b^a R$ from the definition.

The Cauchy Index (continued)

- **Preliminaries**

- A : a rectangle defined by $[a_1, a_2] \times [a_3, a_4]$ which encloses a zero.
- $F = (p, q)$ does not vanish on the boundary of A , ∂A .
- $p \cdot q$ is not zero at each vertex of A .
- Let

$$R_1 = \frac{q(a_1, y)}{p(a_1, y)}, R_2 = \frac{q(a_2, y)}{p(a_2, y)}, R_3 = \frac{q(x, a_3)}{p(x, a_3)}, R_4 = \frac{q(x, a_4)}{p(x, a_4)}.$$

Then, we set (for counterclockwise traversal of ∂A)

$$I_A F = I_{a_4}^{a_3} R_1 + I_{a_3}^{a_4} R_2 + I_{a_1}^{a_2} R_3 + I_{a_2}^{a_1} R_4.$$

- **Proposition***

•T. Sakkalis, "The Euclidean Algorithm and the Degree of the Gauss Map", SIAM J. Computing. Vol. 19, No. 3, 1990.

$I_A F$ is an even integer and the multiplicity $d = -\frac{1}{2} I_A F$.

Illustrative Example for Multiplicity Computation Using the Cauchy Index

- $p(x) = (x-1/2)^5 = 0$
- A root of $p(x)$, $[a] = [0.49, 0.51]$.
- $P(z)$; ($z = x+iy$)

$$p(z) = \left(x + iy - \frac{1}{2}\right)^5 = f(x, y) + ig(x, y)$$

- Create

$$A = [0.49, 0.51] \times [-0.01, 0.01], \quad a_1 = 0.49, a_2 = 0.51, a_3 = -0.01, a_4 = 0.01$$

- Calculate the Cauchy index
 - Roots of $f(x, a_3) = 0$
 - Calculation of

$$I_{a_1}^{a_2} R_3 = -3$$

No.	Roots of $f(x, a_3) = 0$ in $[0, 1]$ (from the IPP)
1	[0.46922316412099, 0.46922316512099]
2	[0.49273457408967, 0.492734576204823]
3	[0.499999997363532, 0.500000001889623]
4	[0.507265424645288, 0.507265426808589]
5	[0.530776834861365, 0.530776835861365]

- Roots No. 2, 3, and 4 are selected since they lie within the interval $[a]$.

Illustrative Example (Continued)

- Similarly, $I_{a_3}^{a_4} R_2 = -2$, $I_{a_2}^{a_1} R_4 = 3$, $I_{a_4}^{a_3} R_1 = 2$
- Calculate $I_A F = I_{a_4}^{a_3} R_1 + I_{a_3}^{a_4} R_2 + I_{a_1}^{a_2} R_3 + I_{a_2}^{a_1} R_4 = -10$
- The multiplicity m of the root is $d = -\frac{1}{2} I_A F = 5$

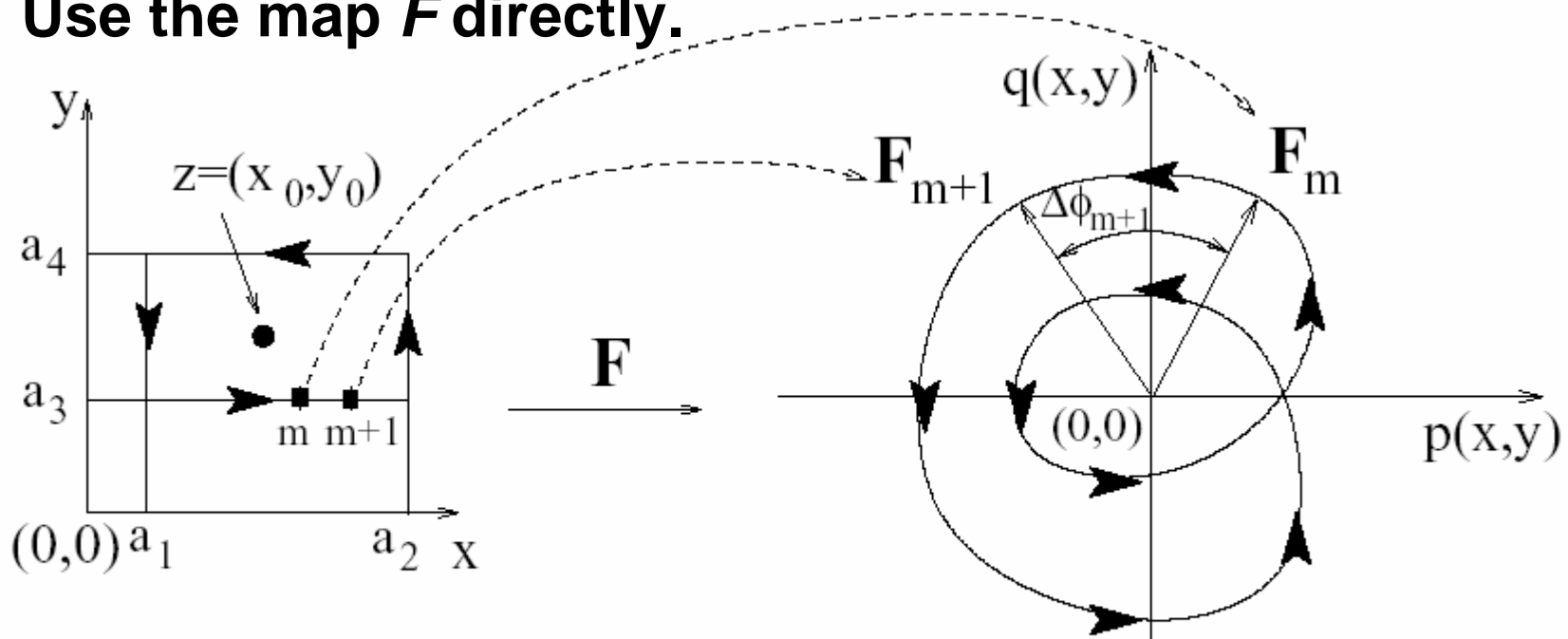
Note

– $I_a^b R = -I_b^a R$.

– *Counterclockwise orientation of ∂A is assumed.*

Direct Computation Method

- Use the map F directly.

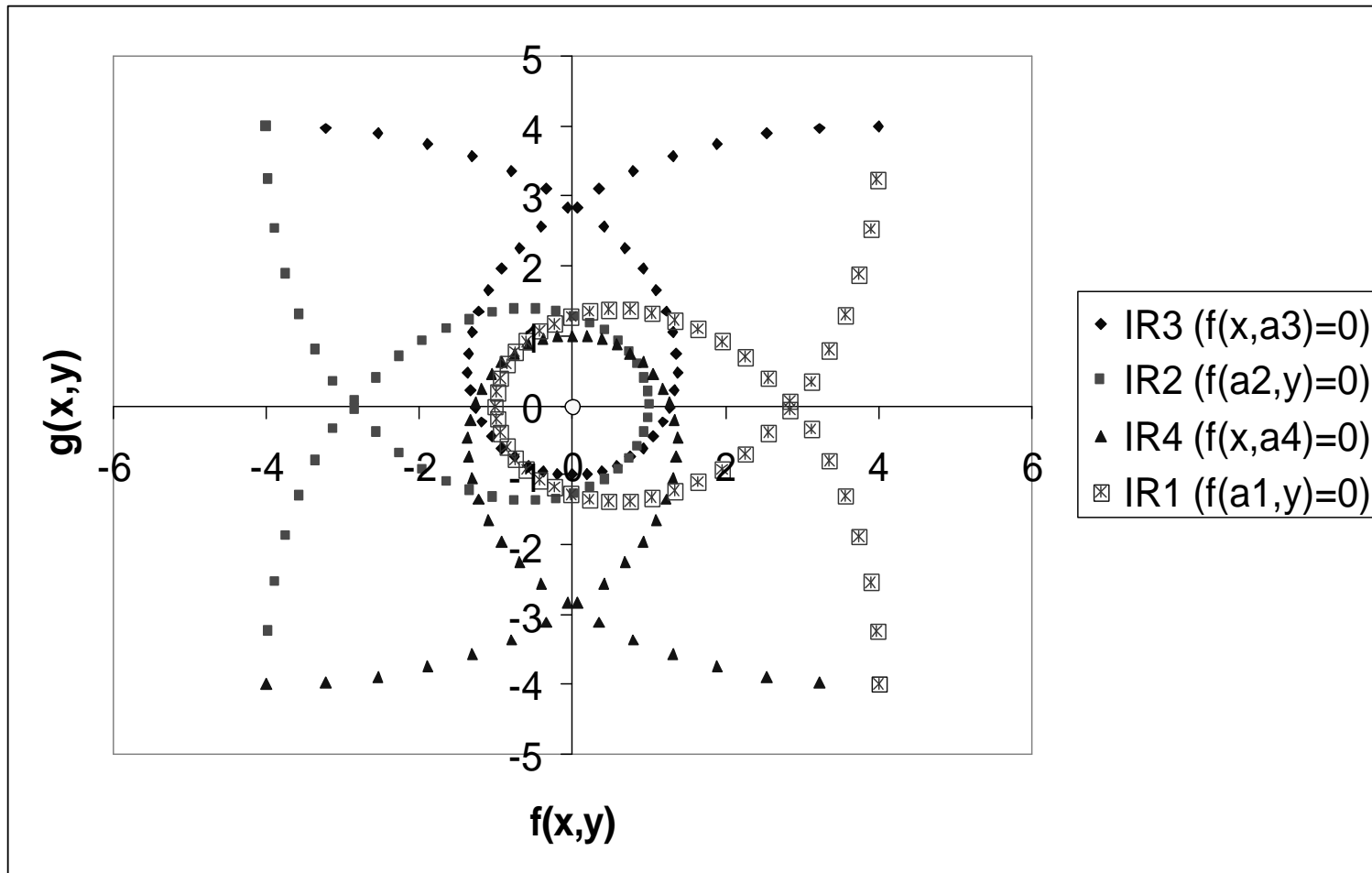


$$\phi_{total} = \sum_{i=0}^n \Delta\phi_{i+1}$$

$$d = \frac{\phi_{total}}{2\pi}$$

Direct Computation Method

$$F : \mathbf{R}^2 \rightarrow \mathbf{R}^2, F(x, y) = (f(x, y), g(x, y)). \quad G : \partial A \rightarrow S^1, G = F / \|F\|,$$



Bisection Algorithm for Solving Univariate Polynomial Equations

- **Univariate polynomial in complex variable z .**
(Substitute x with a complex variable $z = x+iy$)

$$p(z) = a_n z^n + a_{n-1} z^{n-1} + \dots + a_1 z^1 + a_0 = 0$$

- **Input :**

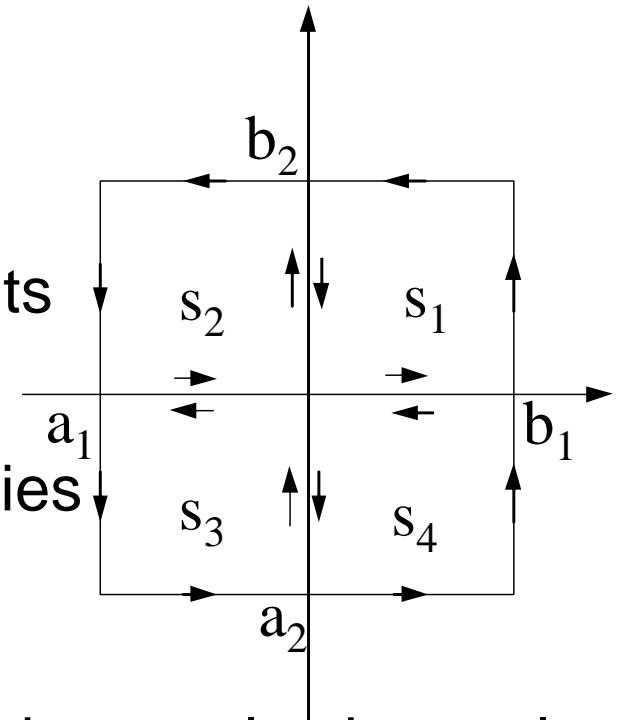
- initial domain : $S = [a_1, b_1] \times [a_2, b_2]$
- a complex polynomial : $p(z)$
- tolerance, number of sample points

- **Output**

- real and complex roots, multiplicities

- **Algorithm**

- Quadtree decomposition
- Direct degree computation method : complex interval arithmetic.



Examples

- Wilkinson polynomial

$$p(t) = \prod_{i=1}^{20} \left(t - \frac{i}{20} \right)$$

No.	Multiplicity	Roots
1	1	[0.05,0.05]+i[-5.769e-10,5.769e-10]
2	1	[0.1,0.1]+i[-5.866e-10,5.866e-10]
3	1	[0.15,0.15]+i[-5.947e-10,5.947e-10]
4	1	[0.2,0.2]+i[-5.947e-10,5.947e-10]
5	1	[0.25,0.25]+i[-5.898e-10,5.898e-10]
6	1	[0.3,0.3]+i[-5.792e-10,5.792e-10]
7	1	[0.35,0.35]+i[-5.792e-10,5.792e-10]
8	1	[0.4,0.4]+i[-5.792e-10,5.792e-10]
9	1	[0.45,0.45]+i[-5.792e-10,5.792e-10]
10	1	[0.5,0.5]+i[-5.745e-10,5.745e-10]
11	1	[0.55,0.55]+i[-5.745e-10,5.745e-10]
12	1	[0.6,0.6]+i[-5.745e-10,5.745e-10]
13	1	[0.65,0.65]+i[-5.745e-10,5.745e-10]
14	1	[0.7,0.7]+i[-5.745e-10,5.745e-10]
15	1	[0.75,0.75]+i[-5.745e-10,5.745e-10]
16	1	[0.8,0.8]+i[-5.745e-10,5.745e-10]
17	1	[0.85,0.85]+i[-5.745e-10,5.745e-10]
18	1	[0.9,0.9]+i[-5.745e-10,5.745e-10]
19	1	[0.95,0.95]+i[-5.745e-10,5.745e-10]
20	1	[1,1]+i[-5.747e-10,5.747e-10]

- Complicated Polynomial (degree 22)

$$p(t) = (t^2 + t + 1)^2 (t - 1)^4$$

$$(t^3 + t^2 + t + 1)^3 (t - 2)(t - 4)^4$$

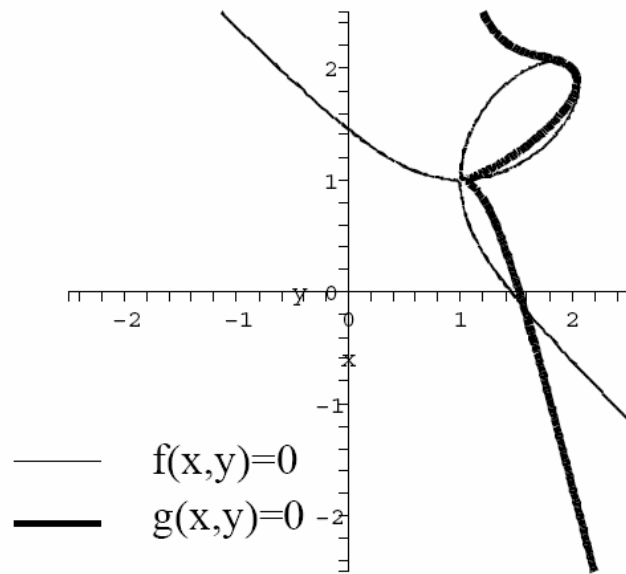
No.	Multiplicity	Roots
1	3	[-5.956e-10,5.956e-10]+i[1,1]
2	4	[1,1]+i[-5.956e-10,5.956e-10]
3	4	[4,4]+i[-5.939e-10,5.939e-10]
4	2	[-0.5,-0.5]+i[0.866,0.866]
5	3	[-1,-1]+i[-5.956e-10,5.956e-10]
6	2	[-0.5,-0.5]+i[-0.866,-0.866]
7	3	[-5.956e-10,5.956e-10]+i[-1,-1]
8	1	[2,2]+i[-5.94e-10,0]

Solving a Bivariate Polynomial System

- **Change of Coordinates**
 - CR : f and g are regular in y .
 - CU : whenever two points (x_0, y_0) and (x_1, y_1) satisfy $f=g=0$, then $y_0=y_1$.
- **Solving a Bivariate Polynomial System**
 - Let f, g satisfy CR and CU and let $h(x)=\text{Res}_y(f, g)$. Then the roots of the system $f=g=0$ are in a one to one correspondence with the roots of $h(x)$. Moreover, $z_i=(x_i, y_i)$ is a real root if and only if x_i is a real root of $h(x)$.
 - Let $h(x)=\text{Res}_y(f, g)$ and $l(y)=\text{Res}_x(f, g)$ and $a_{ij}=[t_i, t_{i+1}]x[s_j, s_{j+1}]$ where in each subinterval $[t_i, t_{i+1}]$ or $[s_j, s_{j+1}]$ there exist precisely one root of $h(x)$ and $l(y)$, respectively. If a_{ij} encloses a real root of $f=g=0$, then the following must be true

$$0 \in f([t_i, t_{i+1}], [s_j, s_{j+1}]) \times g([t_i, t_{i+1}], [s_j, s_{j+1}])$$

Solving a Bivariate Polynomial System : Example



$$\begin{aligned}
 f(x, y) &= x^3 - 3x^2 + 5x - 4 + y^3 \\
 &\quad - 3y^2 + 5y - 2xy = 0, \\
 g(x, y) &= 2x^3 - 2x^2 + x - 4 - 4x^2y + 2xy \\
 &\quad + 9y + 3xy^2 - 8y^2 + y^3 = 0,
 \end{aligned}$$

$$\begin{aligned}
 h(x) &= 56x^9 - 704x^8 + 3880x^7 - 12304x^6 \\
 &\quad + 24744x^5 - 32736x^4 + 28504x^3 \\
 &\quad - 15760x^2 + 5024x - 704.
 \end{aligned}$$

$$\begin{aligned}
 l(y) &= -56y^9 + 608y^8 - 2824y^7 + 7312y^6 \\
 &\quad - 11496y^5 + 11136y^4 - 6328y^3 \\
 &\quad + 1744y^2 - 32y - 64.
 \end{aligned}$$

Root (x,y)	d
$[0.999999978, 1.000000001] \times [0.999999994, 1.000000001]$	5
$[1.57142855, 1.57142859] \times [-0.142857209, -0.142857134]$	1
$[1.99999999, 2.000000003] \times [1.99999996, 2.000000003]$	3

Elimination Theory

I. Resultants

- Sylvester Resultant
- Macaulay Resultant
- Sparse Resultant
- D-Resultant

II. Groebner Bases

III. Symbolic System Solving

Elements of Resultant Theory

- Let:

$$a(t) = a_n t^n + \cdots + a_1 t + a_0$$

$$b(t) = b_m t^m + \cdots + b_1 t + b_0$$

- non zero polynomials, with complex coefficients.
- The resultant of a, b wrt t (or the t -resultant), $\text{Res}_t(a, b) = R$ is

$$R = \begin{vmatrix} a_n & a_{n-1} & \cdots & a_0 & & & \\ & a_n & \cdots & a_1 & a_0 & & \\ & & \ddots & & & \ddots & \\ & & & a_n & \cdots & \cdots & a_0 \\ b_m & \cdots & \cdots & \cdots & b_0 & & \\ & \ddots & & & & \ddots & \\ & & b_m & \cdots & \cdots & \cdots & b_0 \end{vmatrix}$$

- Observe that $\text{Res}_t(a, b) \in \mathbb{C}$.

Properties of the Resultant

- Let us see some well known properties of the resultant:
- **Property 1.** There exist polynomials $A(t), B(t) \in C[t]$ of degrees respectively, $n' < m, m' < n$ so that

$$a(t)A(t) + b(t)B(t) = \text{Res}_t(a, b). \quad (1)$$

- **Property 2.** $\text{Res}_t(a, b) = 0 \iff a(t)$ and $b(t)$ have a common factor of positive degree.

- **Property 3.**

- Let ,
$$a(x, y) = a_n y^n + a_{n-1}(x) y^{n-1} + \dots + a_0(x)$$

$$b(x, y) = \sum_{i=0}^m b_{m-i}(x) y^{m-i} \in k[y][x]$$

- with a_n or $b_m \in C^*$, and consider $p(x) = \text{Res}_y(a, b)$. If x_0 is a root of $p(y)$, then there exists $y_0 \in C$ with the property $a(x_0, y_0) = b(x_0, y_0) = 0$

Cramer's Rule

- Let $f(x,y), g(x,y) \in C[x,y]$ two nonconstant polynomials, and let u, v be indeterminates.

- Consider

$$F : C^2 \rightarrow C^2, \quad F = (f, g)$$

$$A(x, u, v) = \text{Res}_y(f - u, g - v),$$

$$B(y, u, v) = \text{Res}_x(f - u, g - v)$$

- with $F(0,0) = (0,0)$.

Cramer's Rule

- **Theorem**[Cramer's Rule] F has a polynomial inverse if and only if:

$$A(x, u, v) = ax + A_0(u, v),$$

- and

$$B(y, u, v) = by + B_0(u, v), \quad \text{with } ab \neq 0$$

- Moreover, if

$$G(x, y) := \left(-\frac{A_0(x, y)}{a}, -\frac{B_0(x, y)}{b} \right),$$

- Then G is the inverse of F .
- In addition,

$$\deg F = \deg F^{-1}$$